

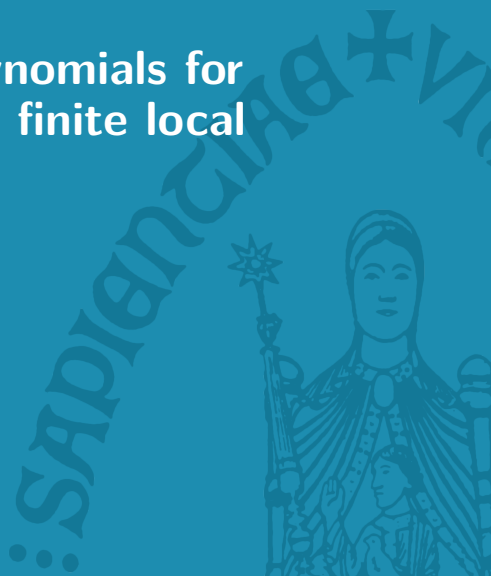
Multiplication polynomials for elliptic curves over finite local rings

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Joint work with Daniele Taufer

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0 Outline

- ① Introduction
- ② Our setting
- ③ Multiplication polynomials
- ④ Consequences

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1 Elliptic curves

Definition

An elliptic curve E is the set of points $(X : Y : Z) \in \mathbb{P}^2(\mathbb{K})$ satisfying a Weierstrass equation, i.e.

$$y^2z + a_1xyz + a_3yz^2 = x^3 + a_2x^2z + a_4xz^2 + a_6z^3$$

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Remark

If $\text{char}(\mathbb{K}) \notin \{2, 3\}$ we can work with the short Weierstrass equation

$$y^2z = x^3 + Axz^2 + Bz^3,$$

without loss of generality.

1 The group structure

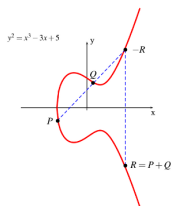
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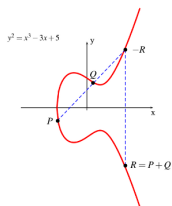
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Theorem

Let p be a prime number, and E an elliptic curve defined over \mathbb{F}_p . There are positive integers $n, k \in \mathbb{Z}_{\geq 1}$ such that $n|(p-1)$ and

$$E(\mathbb{F}_p) \cong \mathbb{Z}/n\mathbb{Z} \times \mathbb{Z}/nk\mathbb{Z}.$$

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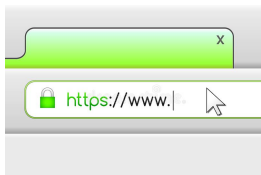
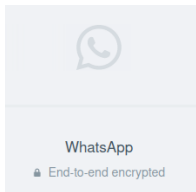
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- ▶ many modern cryptosystems (including *WhatsApp* and *TLS*) are based on ECDLP.



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- ▶ the third case (i.e. $X \in R_k^*, Y, Z \notin R_k^*$) cannot happen because of the Weierstrass equation.

2 Subgroup at infinity

- ▶ We also define a projection $\pi : R_k \xrightarrow{\text{mod } \epsilon} \mathbb{F}_p$ sending

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- ▶ we restrict our attention to E^∞ .

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Proposition

There exist a polynomial $f \in R_k[x]$ such that $x^3 | f$, and for every point $P = (X : 1 : Z) \in E^\infty$ it holds $Z = f(X)$.

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Corollary

Given $P = (P_x : 1 : P_z)$, $Q = (Q_x : 1 : Q_z) \in E^\infty$ it holds

$$(P + Q)_x \in \langle P_x, Q_x \rangle.$$

In particular, $(nP)_x \in \langle P_x \rangle$.

3 Multiplication polynomials

As a consequence, for $P = (X : 1 : f(X))$ we can write

$$(nP)_x = \sum_{i=1}^{k-1} \psi_i(n) X^i.$$

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Remark

ψ_i is well defined as a function; it is not clear yet that this should be a polynomial in n .

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Remark

It can be shown directly that $\psi_1(n) = n$, $\psi_2(n) = \binom{n}{2}a_1 = \frac{n(n-1)}{2}a_1$.

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Theorem (I. and Taufer, 2023)

$\psi_i(n)$ is a degree- i polynomial in $\mathbb{Q}[a_1, \dots, a_6][n]$ with no constant term. Moreover, no primes greater than i appears in the denominators of $\psi_i(n)$.

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Proposition

Let $E(R_k)$ be a curve, $P \in E$. It holds

$$(pP)_x \equiv \psi_p(p)X^p \pmod{X^{p+1}}.$$

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- ▶ if $\nu(P) \neq \nu(Q)$, $\nu(P + Q) = \min(\nu(P), \nu(Q))$;

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- ▶ if $p \nmid n$, $\nu(nP) = \nu(P)$;
- ▶ $\nu(pP) = p\nu(P)$ (assuming $\psi_p(p) \in R_k^*$).

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Lemma

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Theorem (I. and Taufer, 2023)

Let E be an elliptic curve over R_k s.t. $\#E(\mathbb{F}_p) \neq p$ and $\psi_p(p) \in R_k^*$.
Then

$$E \cong E(\mathbb{F}_p) \times \prod_{\substack{1 \leq m \leq k-1 \\ (m,p)=1}} \mathbb{Z}/p^{l_m}\mathbb{Z}, \text{ where } l_m = \left\lfloor \log_p \frac{k-1}{m} \right\rfloor + 1.$$

4 The ECDLP

Let

$$P_x = a_1\epsilon + a_2\epsilon^2 + \dots + a_{k-1}\epsilon^{k-1}$$

and

$$n = b_0 + b_1p + b_2p^2 + \dots + b_dp^d.$$

Suppose we know P , $Q = nP$ and want to recover n . Then:

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$$(Q - b_0P)_x \equiv a'_1b_1\epsilon^p \pmod{\epsilon^{p+1}}$$

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$$(Q - b_0P)_x \equiv a'_1 b_1 \epsilon^p \pmod{\epsilon^{p+1}} \Rightarrow b_1 = (Q - b_0P)_x^{(p)} \cdot (a'_1)^{-1}$$

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Theorem (I. and Tauber, 2023)

It holds

$$b_i = \left(\left(Q - \sum_{j=1}^{i-1} b_j p^j P \right)_x \bmod \epsilon^{m_i+1} \right) / \left((p^i P)_x \bmod \epsilon^{m_i+1} \right),$$

where $m_i = \nu(p^i P)$. Over E^∞ , the discrete logarithm can hence be solved in time $\mathcal{O}(\log(p) \log(n))$. As a consequence, the discrete logarithm over $E(R_k)$ can be efficiently reduced to the corresponding logarithm over $E(\mathbb{F}_p)$.

Thank you for your attention.